

Prise en main du système Hello World

MI074 - 2

Objectif

1/2

L'objectif de cette séance est de présenter en détails la programmation de la première application sur la plateforme. Nous allons voir ainsi (pas forcément dans cet ordre) :

- le mapping mémoire
- le code de boot
- le makefile
- la synchronisation entre les processeurs
- la commande de périphérique
- les directives au compilateur en C
- l'allocation des piles de démarrage

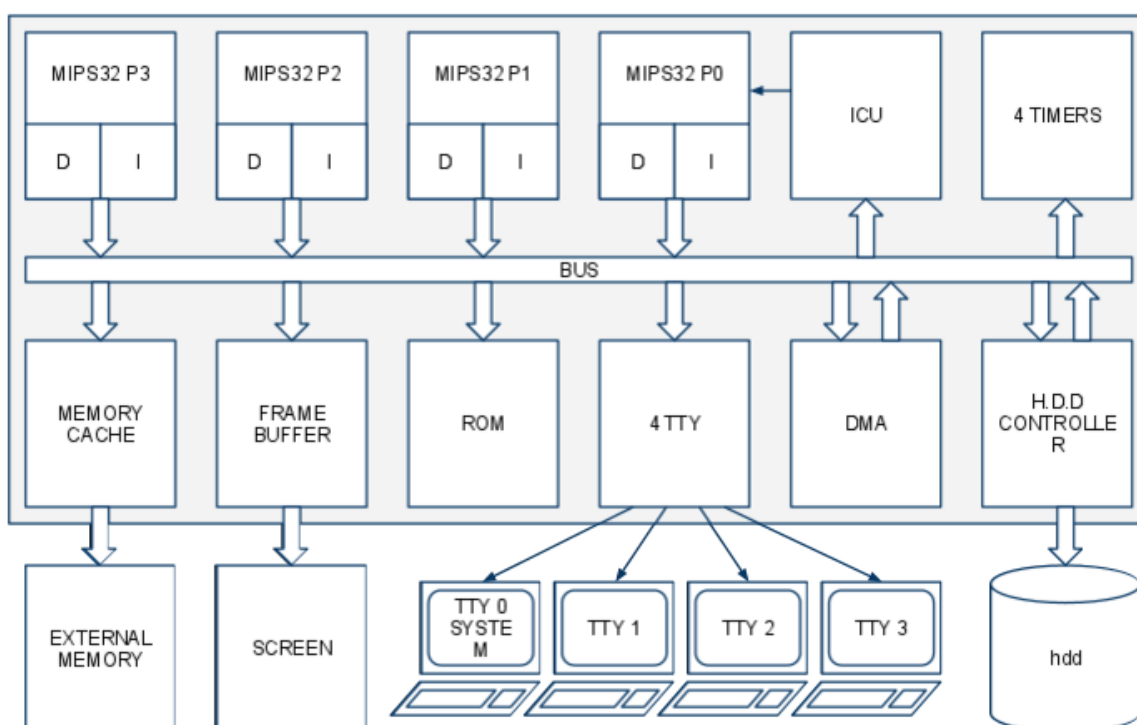
Ce séance est une préparation directe de la séance de TME

Les programmes supports ne seront pas dans l'OS.
C'est seulement préparatoire.

Chaque processeur affiche "Hello World" sur "son" TTY

- entièrement en assembleur
- en assembleur et en C

La plateforme matérielle



description de la mémoire

fichier : segmentation.h

```
// ----- Device mapped segments

#define TIMER_BASE 0xd3200000
#define TIMER_SIZE 0x00000080

#define ICU_BASE 0xd2200000
#define ICU_SIZE 0x00000020

#define DMA_BASE 0xd1200000
#define DMA_SIZE 0x00000014

#define TTY_BASE 0xd0200000
#define TTY_SIZE 0x00000040

#define BD_BASE 0xd5200000
#define BD_SIZE 0x20

#define FB_XSIZE 512
#define FB_YSIZE 512
#define FB_BASE 0x52200000
#define FB_SIZE FB_XSIZE*FB_YSIZE*2
```

```
// ----- ROM mapped segments

#define KTEXT_LMA_BASE 0xbf800000
#define KTEXT_LMA_SIZE 0x00020000

#define KDATA_LMA_BASE 0xbf820000
#define KDATA_LMA_SIZE 0x00020000

#define BOOT_BASE 0xbfc00000
#define BOOT_SIZE 0x00001000

// ----- RAM

#define RAM_BASE 0x7F400000
#define RAM_SIZE 0x01000000

// ----- Application mapped segments

#define KTEXT_BASE 0x80000000
#define KDATA_BASE 0x80020000
#define KDATA_SIZE 0x003E0000

#define USR_TEXT_BASE RAM_BASE
#define USR_TEXT_SIZE 0x00060000
#define USR_DATA_BASE \
    USR_TEXT_BASE + USR_TEXT_SIZE
#define USR_DATA_SIZE 0x00B9F000
```

Description des périphériques

<https://www.soclib.fr/trac/dev/wiki/Component>

pour le TTY :

<https://www.soclib.fr/trac/dev/wiki/Component/VciMultiTty>



wiki: Component / VciMultiTty

[SocLib Components General Index](#)

VciMultiTty

1) Functional Description

This VCI target is a TTY terminal controller. This hardware component controls one or several independent terminals. The number of emulated terminals is defined by the arguments in the constructor (one name per terminal).

Each terminal is acting both as a character display, and a keyboard interface. For each terminal, a specific IRQ is activated when a character entered at the keyboard is available in a buffer. IRQ is kept low as long as the buffer is not empty.

This hardware component checks for segmentation violation, and can be used as a default target.

This component uses a [TtyWrapper](#) per terminal in order to abstract the simulated ttys. The terminal index *i* is defined by the ADDRESS[12:4] bits.

Each TTY controller contains 3 memory mapped registers:

- TTY_WRITE

This 8 bits pseudo-register is write only. Any write request will interpret the 8 LSB bits of the WDATA field as an ASCII character, and this character will be displayed on the addressed terminal.

- TTY_STATUS

This Boolean status register is read-only. A read request returns the zero value if there is no pending character. It returns a non zero value if there is a pending character in the keyboard buffer.

- TTY_READ

Description des périphériques

Sur la page de description du périphérique on trouve :

The file `source:trunk/soclib/soclib/module/connectivity_component/vci_multi_tty/include/soclib/tty.h` defines `TTY_WRITE`, `TTY_STATUS`, `TTY_READ` and `TTY_SPAN`.

fichier : `devices.n`

```
#ifndef _DEVICES_H_
#define _DEVICES_H_

/* TTY mapped registers offset */
#define TTY_SPAN 4
#define TTY_WRITE_REG 0
#define TTY_STATUS_REG 1
#define TTY_READ_REG 2

#endif
```

Description de la mémoire pour le linker

fichier `kldscript.h`

```
#include "segmentation.h"

MEMORY
{
    boot : ORIGIN = BOOT_BASE, LENGTH = BOOT_SIZE
}

SECTIONS
{
    .boot : { *(.boot) } > boot
}
```

Description de la mémoire pour le linker

fichier kldscript.h

```
#include "segmentation.h"
MEMORY
{
    boot : ORIGIN = BOOT_BASE, LENGTH = BOOT_SIZE
}
SECTIONS
{
    .boot : { *(.boot) } > boot
}

#include "segmentation.h"
MEMORY
{
    boot : ORIGIN = BOOT_BASE, LENGTH = BOOT_SIZE
}
SECTIONS
{
    .boot : { *(.boot) *(.text) *(.boot.*) *(.data) *(.rodata*) } > boot
}
```

le code de boot

fichier : boot.S

```
#include <devices.h>
#include <segmentation.h>

.section .boot, "ax", @progbits
.ent boot
.align 2

boot:
    li    $26, 0
    mtc0  $26, $12    # Status Register
    mfc0  $4, $0
    la    $5, str
    j     tty_puts

// Function name : fputs
// Arguments    : $4 <-> processor id, $5 <-> @str
// Description  : print charecters of string str
// -----
tty_puts:
    la    $26, TTY_BASE
    li    $27, TTY_SPAN * 4
    multu $4, $27
    mflo  $27
    addu  $26, $26, $27
loop:
    lbu   $27, 0($5)
    beqz  $27, deadLoop
    sw    $27, (TTY_WRITE_REG)($26)
    addiu $5, $5, 1
    j     loop
deadLoop:
    j     deadLoop

str: .asciiz "hello world"

.end boot
```

Signification des drapeaux

Section Flag Characters	
Flag Characters	Description
w	Write access allowed.
a	Section is allocated in memory.
x	Section contains executable instructions.
s	Section contains "short" data.
o	Section adds ordering requirement. The 'o' flag is only for <code>ELF</code> (Unix*) files.

Section Types	
Section Type	Description
"progbits"	Sections with initialized data or code.
"nobits"	Sections with uninitialized data (<code>bss</code>).
"comdat"	<code>COMDAT</code> sections, Windows NT specific. See Windows NT (COFF32) Specific Section Flag Operands .
"note"	Note sections.

registre status

dans MIPS_vol3 p. 53

contenu des 16 bits de poids faible du registre **SR**:

IM[7:0]	0	0	0	UM	0	ERL	EXL	IE
---------	---	---	---	----	---	-----	-----	----

Cette version du processeur MIP32 n'utilise que 12 bits du registre SR :

IE : Interrupt Enable

EXL : Exception Level

ERL : Reset Level

UM : User Mode

IM[7:0] : Masques individuels pour les six ligne d'interruption matérielles (bits IM[7:2]) et pour les 2 interruptions logicielles (bits IM[1:0])

registre status

IM[7:0]	0	0	0	UM	0	ERL	EXL	IE
---------	---	---	---	----	---	-----	-----	----

- Le processeur a le droit d'accéder aux ressources protégées (registres du CP0, et adresses mémoires > 0x7FFFFFFF) si et seulement si le bit UM vaut 0, ou si l'un des deux bits ERL et EXL vaut 1.
- Les interruptions sont autorisées si et seulement si le bit IE vaut 1, et si les deux bits ERL et EXL valent 00, et si le bit correspondant de IM vaut 1.
- Les trois types d'événements qui déclenchent le branchement au GIET (interruptions, exceptions et appels système) forcent le bit EXL à 1, ce qui masque les interruptions et autorise l'accès aux ressources protégées.
- L'activation du signal RESET qui force le branchement au Boot-Loader force le bit ERL à 1, ce qui masque les interruptions et autorise l'accès aux ressources protégées.
- L'instruction ERET force le bit EXL à 0.

```
INCLUDES = segmentation.h
KERNEL   = kernel-soclib.bin
CSRC     =
SSRC     = boot.S
KOBJS    = $(SSRC:%.S=%o) $(CSRC:%.c=%o)
```

```
.PHONY: clean realclean
```

```
# CC tools and parameters
```

```
#-----
BUILD_DIR := $(shell pwd)
CCTOOLS   ?=
CPU       = mipsel
CC        = $(CCTOOLS)/bin/$(CPU)-unknown-elf-gcc
LD        = $(CCTOOLS)/bin/$(CPU)-unknown-elf-ld
OD        = $(CCTOOLS)/bin/$(CPU)-unknown-elf-objdump
SIMUL     = ../bin/simulation.x
```

```
CFLAGS    = -fno-builtin -l. -fomit-frame-pointer -O3 -G0 -Wall -Werror -mips32r2
LDFLAGS   =
TRASH     = /dev/null||true
```

```
# Kernel building rules
```

```
#-----
all: depend $(KERNEL)
```

Makefile

Makefile

```
$(KERNEL): $(KOBJS) kldscript
    $(LD) -o $@ $(KOBJS) $(LDFLAGS) -T$(BUILD_DIR)/kldscript
    $(OD) $@ -D > $@.dump

kldscript : kldscript.h segmentation.h
    cpp $< | egrep -v "#|/" | grep . > $@

clean:
    $(RM) vcitty* *.bak *.o *.pdf *~ kldscript *.dump $(KERNEL) 2> $(TRASH)

simul: $(KERNEL)
    $(SIMUL)

pdf:
    a2ps -1 --medium=A4 --file-align=fill -o - -l100 \
        Readme.txt Makefile *.h $(SSRC) $(CSRC) |\
    ps2pdf -sPAPERSIZE=a4 - `basename $$PWD`.pdf

#-----
%.o: %.c
    $(CC) $(CFLAGS) -c $< -o $@
%.o: %.S
    $(CC) $(CFLAGS) -c $< -o $@

depend: ; makedepend -i. $(CSRC)
```

Compilation & Exécution

```
> make
makedepend -i.
/usr/local/cctools/bin/mipsel-unknown-elf-gcc -fno-builtin -i. -fomit-frame-pointer -O3 -G0 -
Wall -Werror -mips32r2 -c boot.S -o boot.o
cpp kldscript.h | egrep -v "#|/" | grep . > kldscript
/usr/local/cctools/bin/mipsel-unknown-elf-ld -o kernel-soclib.bin boot.o -
T/Users/franck/Enseig/ue1_os/ose/tp1/e0/kldscript
/usr/local/cctools/bin/mipsel-unknown-elf-objdump kernel-soclib.bin -D > kernel-soclib.bin.
dump
```

```
> make simul
../../bin/simulation.x
```

```
SystemC 2.2.0 --- Nov 9 2009 17:32:16
Copyright (c) 1996-2006 by all Contributors
ALL RIGHTS RESERVED
```

```
Mapping table: ad:(8) id:(8) cacheability mask: 0xf000000
<Segment "boot": @0x0xbfc00000, 0x0x1000 bytes, (0),
<Segment "ktext": @0x0xbf800000, 0x0x20000 bytes, (0),
<Segment "kdata": @0x0xbf820000, 0x0x20000 bytes, (0),
....
```



Ajout de code C / boot

La différence est qu'il va falloir réserver une pile pour chaque processeur

```
#include <segmentation.h>
#include <config.h>

#define STACK_SIZE CONFIG_BOOT_STACK_SIZE

.section .boot,"ax",@progbits
.extern __do_init
.ent boot
.align 2

boot:
    li    $26, 0
    mtc0  $26, $12           # Status Register
    mfc0  $4,  $0           # CPU_ID
    la    $27, RAM_BASE+RAM_SIZE # top memory
    li    $26, (STACK_SIZE)
    mult  $4, $26
    mflo  $29
    subu  $29, $29, $27     # $29 <= TOP - procid * STACK_SIZE
    addiu $29, $29, -1*4   # for __do_init argument
    la    $26, __do_init
    jr    $26

.end boot
```

Ajout de code C / __do_init

que se passerait-il si on n'utilisait pas de pile séparée ?

```
#include <segmentation.h>
#include <config.h>
#include <devices.h>

void tty_puts(unsigned tty_id, char *buffer)
{
    unsigned volatile *tty_base = (unsigned *) TTY_BASE + (tty_id * TTY_SPAN);
    while (*buffer)
        tty_base[TTY_WRITE_REG] = *buffer++;
}

void __do_init(unsigned cpu_id)
{
    tty_puts(cpu_id, "Hello World!\n");
    while (1);
}
```

Ajout de code C / kldscript

code objet issu du code C

Disassembly of section .text:

```
00000000 <tty_puts>:  
0: 80a20000 lb v0,0(a1)  
4: 00042100 sll a0,a0,0x4  
8: 3c03d020 lui v1,0xd020  
c: 10400006 beqz v0,28 <tty_puts+0x28>
```

```
...  
28: 03e00008 jr ra  
2c: 00000000 nop
```

```
00000030 <__do_init>:  
30: 3c02d020 lui v0,0xd020  
34: 00042100 sll a0,a0,0x4  
38: 3c030000 lui v1,0x0
```

```
...  
5c: 08000017 j 5c <__do_init+0x2c>  
60: 00000000 nop
```

Disassembly of section .reginfo:

```
00000000 <.reginfo>:  
0: 8000003c lb zero,60(zero)  
...
```

Disassembly of section .pdr:

```
00000000 <.pdr>:  
...  
18: 0000001d 0x1d  
1c: 0000001f 0x1f  
...  
38: 0000001d 0x1d  
3c: 0000001f 0x1f
```

Disassembly of section .rodata.str1.4:

```
00000000 <$.LC0>:  
0: 6c6c6548 0x6c6c6548  
4: 6f57206f 0x6f57206f  
8: 21646c72 addi a0,t3,27762  
c: 0000000a movz zero,zero,zero
```

Ajout de code C / kldscript

que se passerait-il si on n'utilisait pas de pile séparée ?

```
#include "segmentation.h"  
  
MEMORY  
{  
    boot : ORIGIN = BOOT_BASE, LENGTH = BOOT_SIZE  
}  
  
SECTIONS  
{  
    .boot : { *(.boot) *(.text) *(.boot.*) *(.data) *(.rodata*) } > boot  
}
```